

Firewall Assignment

Assignment 8

Network Security (CS6903)

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Abstract

We have created a firewall system. A firewall monitors the incoming and outgoing traffic in the network A firewall is a network security device that monitors incoming and outgoing network traffic and permits or blocks data packets based on a set of security rules.

First, we try to implement a simple firewall system with two network interface cards connecting to the external network (Internet) and the internal network which is supposed to be secured. In this system we hardcore a simple rules set. Then we improve the firewall to make it more advanced by extending the supported rule set up to layer 4 (including MAC, IPv4 IPv6, ICMP for IPv4/v6, TCP/UDP) and not hard-coding the ruleset. Then we analyzed the performance of the implemented firewall, calculating the packet per second the implementation can handle for different scenarios. We have then tried to improve the performance of the firewall. In the final step, we have shown how to detect attacks such as DoS.



Setup

Three VM has been set up in which VM1 is the internal network, VM2 is the firewall and VM3 is the external network that is connected to the internet. We have done the setup on 2 laptops.



Fig: Setup on laptop 2



Configuration

The following are the Ip addresses and the route table configurations of the VM's

Laptop 1

root@vm	2-Standard-PC-Q35-ICH9-2009:/home/vm2# ifcont	fig
enp1s0:	flags=4163 <up,broadcast,running,multicast></up,broadcast,running,multicast>	mtu 1500
	inet 192.168.130.135 netmask 255.255.255.0	broadcast 192.168.130.255
	inet6 fe80::90c1:e57:e845:dbb prefixlen 64	scopeid 0x20 <link/>
	ether 52:54:00:f7:69:35 txqueuelen 1000 (I	Ethernet)
	RX packets 25929 bytes 2081450 (2.0 MB)	
	RX errors 0 dropped 21187 overruns 0 fram	ne O
	TX packets 30200 bytes 2795770 (2.7 MB)	
	TX errors 0 dropped 0 overruns 0 carrier 0	0 collisions 0

Fig: VM1 config



Fig: VM2 config (Firewall)



Fig: VM3 config



estination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface
0.0.0	192.168.130.158	0.0.0.0	UG	0	0	0	enp1s0
0.0.0	192.168.130.1	0.0.0.0	UG	20100	0	0	enp1s0
69.254.0.0	0.0.0	255.255.0.0	U	1000	Θ	Θ	enp1s6
92.168.130.0	0.0.0.0	255.255.255.0	U	100	0	0	enp1s0

Fig: VM1 route tables

Kernel IP rout	ing table						
Destination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface
0.0.0	192.168.130.1	0.0.0.0	UG	100	Θ	Θ	enp1s0
0.0.0	192.168.140.1	0.0.0.0	UG	101	0	0	enp6s0
169.254.0.0	0.0.0.0	255.255.0.0	U	1000	0	Θ	enp1s0
192.168.130.0	0.0.0.0	255.255.255.0	U	100	Θ	Θ	enp1s0
192.168.140.0	0.0.0.0	255.255.255.0	U	101	0	0	enp6s0

Fig: VM2 route table (Firewall)

root@vm3-Standar Kernel IP routir	rd-PC-Q35-ICH9-20 ng table	009:/home/vm3# r	oute -	۱					
Destination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface		
0.0.0.0	192.168.140.1	0.0.0.0	UG	100	Θ	0	enp1s0		
169 254 0 0	0.0.0.0	255.255.0.0		1866	0	0	enp1s8		
192.168.130.0	192.168.140.181	255.255.255.0	UG	0	0	0	enp1s0		
root@vm3-Standar	d-PC-Q35-ICH9-20	009:/home/vm3# i	ptable	s -t na	t-A	POSTROL	JTING -o	enp1s0	-j MASQUERAD

Fig: VM3 route table

Laptop 2



Fig: VM1 config





Fig: VM2 config (Firewall)



Fig: VM 3 config (external)

Testing

To check if the VM's are working properly we are going to ping Google from VM1

Server: Address:	127.0.0.53 127.0.0.53#53	
Non-authoritati Name: google Address: 142.25	ive answer: .com 50.195.206	

Fig: Google ip add lookup



vm1@vm1-VirtualBox:~\$ ping 142.250.195.206
PING 142.250.195.206 (142.250.195.206) 56(84) bytes of data.
64 bytes from 142.250.195.206: icmp_seq=1 ttl=53 time=24.5 ms
64 bytes from 142.250.195.206: icmp_seq=2 ttl=53 time=24.4 ms
64 bytes from 142.250.195.206: icmp_seq=3 ttl=53 time=24.4 ms
64 bytes from 142.250.195.206: icmp_seq=4 ttl=53 time=24.0 ms
64 bytes from 142.250.195.206: icmp_seq=5 ttl=53 time=24.0 ms
^C
142.250.195.206 ping statistics
5 packets transmitted, 5 received, 0% packet loss, time 4006ms
rtt min/avg/max/mdev =_24.030/24.280/24.486/0.205

Fig: ping to google using VM1 reply incoming

il€	e <u>E</u> dit	<u>V</u> iew <u>G</u> o	<u>C</u> apture	<u>A</u> nalyze	<u>S</u> tatisti	ics Telephon <u>y</u>	<u>W</u> irele	ss <u>T</u> ools	<u>H</u> elp			
		ا ۱	0101 0110 0111	X C	٩ (🗭 🏓 警	<u></u>			Ð, e		3
	icmp											
١.	Т	íme	Source		[Destination		Protocol	Length	Info		
	10	.000000000	192.16	8.100.6		142.250.195.2	06	ICMP	100	Echo	(ping)	request
	20	.024473609	142.25	0.195.206		192.168.100.6	j	ICMP	100	Echo	(ping)	reply
	31	.002315061	192.16	8.100.6		142.250.195.2	06	ICMP	100	Echo	(ping)	request
	41	.026709926	142.25	0.195.206		192.168.100.6	j	ICMP	100	Echo	(ping)	reply
	52	.003630023	192.16	8.100.6		142.250.195.2	06	ICMP	100	Echo	(ping)	request
	62	.028024109	142.25	0.195.206		192.168.100.6	j	ICMP	100	Echo	(ping)	reply
	73	.005046845	192.16	8.100.6		142.250.195.2	06	ICMP	100	Echo	(ping)	request
	83	.028976972	142.25	0.195.206	-	192.168.100.6	5	ICMP	100	Echo	(ping)	reply
	94	.005843110	192.16	8.100.6		142.250.195.2	06	ICMP	100	Echo	(ping)	request
	10 4	.029802960	142.25	0.195.206	-	192.168.100.6	j	ICMP	100	Echo	(ping)	reply

Fig: Wireshark capture for above at VM1

File	e Edit View Go (Capture Analyze Statis	stics Telephony Wire	less Tools I	lelp			
1.00	<u></u>		icephon <u>y</u>		<u>-</u> c.p			
	🔳 🙇 🔘 🗎		🗢 🔿 警 有	⊻ 💷 🔳	⊕€		D N	
	icmp							
No.	Time	Source	Destination	Protocol Le	ngth Info			
	109 91.436649136	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	110 91.436682943	192.168.101.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	111 91.460404751	142.250.195.206	192.168.101.6	ICMP	100 Echo	(ping)	reply	id=0
-	112 91.460427795	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply	id=0
	113 92.439220549	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	114 92.439275706	192.168.101.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	115 92.462273797	142.250.195.206	192.168.101.6	ICMP	100 Echo	(ping)	reply	id=0
	116 92.462305956	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply	id=0
	117 93.440569186	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	118 93.440619639	192.168.101.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	119 93.463565456	142.250.195.206	192.168.101.6	ICMP	100 Echo	(ping)	reply	id=0
	120 93.463596568	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply	id=0
	121 94.441851607	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	122 94.441891454	192.168.101.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	123 94.464636223	142.250.195.206	192.168.101.6	ICMP	100 Echo	(ping)	reply	id=0
	124 94.464667083	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply	id=0
	129 95.442855759	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	130 95.442894111	192.168.101.6	142.250.195.206	ICMP	100 Echo	(ping)	request	id=0
	133 95.465297487	142.250.195.206	192.168.101.6	ICMP	100 Echo	(ping)	reply	id=0
	134 95.465329005	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply	id=0

Fig: Wireshark capture for above at firewall



Task 1

Creating a simple Firewall using Socket Programming

In this part, we have implemented a simple firewall. This firewall works at layer 3 i.e IP layer. In this, there is a precoded list of IP addresses that can be blocked. By **default, it allows all the packets**. It checks if the IP address is in the blocked list, if it is there it makes allow = False.



Fig: Hardcoded IP addresses

Command for the simple firewall.



As we can see **142.250.195.206** is the blocked ip address. We have run a ping command from the internal host to **142.250.195.206** i.e to the Google server. The request of the ping request is passed by the firewall but the reply from the Google server is not allowed by the firewall. This happens because the Google IP address is on the blocking list.



Fig: Tried ping command again from internal host VM1



[Ethernet][IPv4][ICMPv4] [Src MAC]: 52:54:00:f7:69:35, [Dstn MAC]: 52:54:00:7f:ab:9d [Src IP]: 192.168.130.135, [Dstn IP]: 142.250.182.14 [Status]: Allowed [PPT] : 0.00015059 [Ethernet][IPv4][ICMPv4] [Src MAC]: 52:54:00:d6:10:87, [Dstn MAC]: 52:54:00:41:10:ec SrcIP]: 142.250.182.14, [Dstn IP]: 192.168.130.135 [Status]: Dropped [PPT] : 0.00015407 [Src MAC]: 52:54:00:d6:10:87, [Dstn MAC]: 52:54:00:41:10:ec [SrcIP]: 142.250.182.14, [Dstn IP]: 192.168.130.135 [Status]: Dropped [PPT] : 0.00014085 [Src MAC]: 52:54:00:41:10:ec, [Dstn MAC]: 52:54:00:d6:10:87 [SrcIP]: 142.250.182.14, [Dstn IP]: 192.168.130.135 [Status]: Dropped [PPT] : 6.409e-05 [Ethernet][IPv4][ICMPv4] [Src MAC]: 52:54:00:f7:69:35, [Dstn MAC]: 52:54:00:7f:ab:9d [Src IP]: 192.168.130.135, [Dstn IP]: 142.250.182.14 [Status]: Allowed [PPT] : 5.913e-05

Fig: Packets allowed and discarded by the blocking process (ScreenShot from Laptop 1)

PCAP analysis:

The PCAP on VM1 shows that there is no reply packet from the Google server (142.250.195.206) since the IP has been blocked by the firewall in the source IP field.



ļ	<u>F</u> ile	<u>E</u> dit	<u>V</u> iew <u>G</u>	o <u>C</u> apture	<u>A</u> nalyze	<u>S</u> tatistics	Telephon	<u>y W</u> ire	less <u>T</u> ools	<u>H</u> elp			
L			۵ ا	8101 8110 8113	XC	۹ (=	• 🔿 🗳		₹		Ð, G		
	icr	mp											
N	о.	1	Time	Source		Des	tination		Protocol	Length	Info		
r	_	3 :	1.15303601	11 192.16	8.100.6	142	2.250.195.	206	ICMP	100	Echo	(ping)	request
T		7 3	2.17944636	52 192.16	8.100.6	142	2.250.195.	206	ICMP	100	Echo	(ping)	request
		8 3	3.20319365	53 192.16	8.100.6	142	2.250.195.	206	ICMP	100	Echo	(ping)	request
		94	4.22736646	53 192.16	8.100.6	142	2.250.195.	206	ICMP	100	Echo	(ping)	request
		10 !	5.25136974	46 192.16	8.100.6	142	2.250.195.	206	ICMP	100	Echo	(ping)	request
		15 (6.27481425	56 192.16	8.100.6	142	2.250.195.	206	ICMP	100	Echo	(ping)	request
		16	7.29900466	68 192.16	8.100.6	142	2.250.195.	206	ICMP	100	Echo	(ping)	request
		17 8	8.32270658	38 192.16	8.100.6	142	.250.195.	206	ICMP	100	Echo	(ping)	request
L	_	18 9	9.34709273	37 192.16	8.100.6	142	2.250.195.	206	ICMP	100	Echo	(ping)	request

Fig: Wireshrk capture at VM1 shows that reply is blocked by firewall

, i	cmp						
lo.	Time	Source	Destination	Protocol	Length Info		
_	28 19.444889897	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	29 19.445041671	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	30 19.468988830	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply
	39 20.471415490	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	40 20.471630814	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	41 20.492873684	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply
	42 21.495020957	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	43 21.495216136	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	44 21.518102391	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply
	45 22.519147480	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	46 22.519355503	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	47 22.542838465	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply
	48 23.543259208	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	49 23.543460414	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	50 23.582578963	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply
	60 24.566550081	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	61 24.566704389	192.168.100.6	142.250.195.206	ICMP	100 Echo	(ping)	request
	62 24.589439880	142.250.195.206	192.168.100.6	ICMP	100 Echo	(ping)	reply
	60 DE E01006046	100 160 100 6	140 050 105 006	TOMD	100 Eabo	(ning)	request

Fig: Reply received at firewall but is not forwarded, it is blocked.



Task 2

Extending the ruleset and its operation on the Firewall

In this task, we need to add additional rules to improve the functionality of the firewall. We are filtering at layer 2 [Ether], layer 3 [IP], and layer 4 [TCP, UDP]. We are doing dynamic rule management by addition, deletion, and updating the rules in the firewall. This is done by storing the rules of the firewall in a json file. The json file is edited in order to update the rules.

Command:

python3 firewall.py

How our Firewall Works



Fig: Firewall CFG



Firewall Home Screen : Home Screen



Fig: Home screen

Blocking Process :

Blocking process looks at every single packet and decides whether it can be allowed or it needs to be blocked based on the rules mentioned in the rules.json. In our Firewall, by **default**, it **blocks all the packets**. If we want to allow any type of packet we need to mention it in rules.json

```
[Ethernet][IPv4][ICMPv4]
[Src MAC]: 52:54:00:f7:69:35, [Dstn MAC]: 52:54:00:7f:ab:9d
[Src IP]: 192.168.130.135, [Dstn IP]: 142.250.182.14
[Status]: Allowed
[PPT] : 0.00015059
[Ethernet][IPv4][ICMPv4]
[Src MAC]: 52:54:00:d6:10:87, [Dstn MAC]: 52:54:00:41:10:ec
[SrcIP]: 142.250.182.14, [Dstn IP]: 192.168.130.135
[Status]: Dropped
[PPT] : 0.00015407
[Ethernet]
[Src MAC]: 52:54:00:d6:10:87, [Dstn MAC]: 52:54:00:41:10:ec
[SrcIP]: 142.250.182.14, [Dstn IP]: 192.168.130.135
[Status]: Dropped
[PPT] : 0.00014085
[Ethernet]
[Src MAC]: 52:54:00:41:10:ec, [Dstn MAC]: 52:54:00:d6:10:87
[SrcIP]: 142.250.182.14, [Dstn IP]: 192.168.130.135
```

Fig: blocking process running

Statistics :

When the Firewall run is complete the firewall program returns the Statistics page which returns the number of allowed and dropped packets. It also returns the average time taken to process each packet.



Firewall Capture Statistics
No of packets allowed : 78
No of packets dropped : 132
Mean Packet Processing Time : 0.0001248
No of rules in system : 3 Press 'c' to continue

Fig: Statistics page

rules.json:

rules.json file is the file from which we read the rules which are used in te blocking process. We can also add, delete or update rules in this.

Ope	n ~ (Fi	rules.json /home/vm1 Save =		×
1				
2	"L2": [
3	{			
4		"rule_id": 168,		
5		"rule": "allow"		
6	}			
7	1,			
8	"L3V4":	[
9	{			
10		"rule_id": 57,		
11		"src_ip": "142.250.182.14",		
12		"rule": "Discard"		
13	},			
14	{			
15		"rule_id": 57,		
16		"src_ip": "192.168.130.135",		
17		"rule": "allow"		
18	}			
19	1,			
20	"L4TCP":			
21	"L4UDP":	[]		
22				
		JSON \checkmark Tab Width: 8 \checkmark Ln 1, Col 1	~	INS

Fig: rules.json file

When "s" is pressed the blocking process is started. When "r" is pressed then the rules. json file is opened in which we can add, remove and update the rules. After every run of the firewall the program prints the statistics of the code.



Types of Rules

L2: Mac Layer Filtering Rule

Filtering packets based on Mac address



Fig: Allowed packet rule

IP Packet Filtering Rule

Filtering packets based on IP rule

```
"L3V4": [
{
    "rule_id": 57,
    "ipv4protocol": 3,
    "rule": "Allow"
}
],
```

Fig: Allowed packet rule

UDP Packet filtering Rule

Filtering packets based on IP rule

```
"L4UDP": [
    {
        "rule_id": 505,
        "udpsrc_port": 403,
        "udpdest_port": 9876,
        "rule": "Allow"
    }
]
```

Fig: Allowed packet rule



TCP Packet filtering Rule

Filtering packets based on IP rule

"L4TCP": [{ "rule id": 903, "tcpsrc_port": 403, "tcpdest port": 5555, "rule": "Allow" }],

Fig: Allowed packet rule



Task 3

Performance examination and improvement

In this part, we have tried to analyze the system. We have used the metric processing time per packet for analyzing the performance of the system. We have increased the number of rules and tried to observe how this affects the performance of the system.

We are generating traffic on VM3 using "generate_traffic.py" which uses the "nping" command. This file generates mixed traffic (i.e TCP/UDP/IP/IP) continuously and sends it to VM1.

Command: python3 traffic_generator.py g

This command will generate a total 1000 packets and send to the destination IP which is hardcoded in the file itself.

At the same time, we also wrote code to create rules using random parameters.

Command : *python3 rules_generator.py rm No_Of_Rules_In_Each_Category*.

No_Of_Rules_In_Each_Category is an integer value stating how many rules will be created for each traffic category like TCP, UDP etc.

**********	******	
STATISTICS	STATISTICS	
******	********	
The Statistics of the system are as follows	The Statistics of the system are as follows	
Average Time Taken to process packet : 7.196308649122883e-05	Average Time Taken to process packet : 8.349324387124151e-05	
No of packets allowed : 1206	No of packets allowed : 1288	
No of packets dropped : 4494	No of packets dropped : 3403	
No of rules in system : 4 Maximum Matching Fields in Rules : 4	No of rules in system : 12 Maximum Matching Fields in Rules : 5	
press enter to continue	press enter to continue	



***********	*******	
STATISTICS		
********	******	
The Statistics of the system	n are as follows	
Average Time Taken to proce	ss packet : 9.80684191347755e-05	
No of packets allowed : 22	10	
No of packets dropped : 380	90	
No of rules in system : 16 Maximum Matching Fields in f	Rules : 5	
***************	**************	
STATISTICS	STATISTICS	
*****	*********	
The Statistics of the system are as follows	The Statistics of the system are as follows	
Average Time Taken to process packet : 0.0001538	Average Time Taken to process packet : 0.0001839954	
No of packets allowed : 2853	No of packets allowed : 3193	
No of packets dropped : 3600	No of packets dropped : 1627	
No of rules in system : 52 Maximum Matching Fields in Rules : 5	No of rules in system : 100 .Maximum Matching Fields in Rules : 5	

Fig: Statistics

Statistics				
No of Rules	Processing time per packet			
4	0.00007196			
8	0.00008349			
16	0.00009806			
50	0.0001538			
100	0.0001839			



Processing Time per Packet

Fig: Analysis of statistics

We can see clearly as the number of Rules is increased the the time taken for processing the packet also increases.



Task 4 - Part B:

Detecting attacks in the network using Firewall

DoS Attack

A Denial-of-Service (DoS) attack is one that attempts to bring a machine or network to a halt, rendering it unreachable to its intended users. DoS attacks work by inundating the target with traffic or delivering it information that causes it to crash. The DoS attack deprives genuine users, such as employees, of the service or resource they expected in both cases.



Fig: DoS Attack (ref: https://www.cloudflare.com/learning/ddos/glossary/denial-of-service/)

Prevention using Firewall

In this task we're going to detect the Dos attack that we are performing on VM1. In order to prevent the DoS attack we have tried to put a cap on the number of packets incoming from one IP. If a lot of IP packets are incoming from one IP then the firewall system detects that there is some DoS attack and then blocks that particular IP. Here is our implementation of the DoS attack:



Fig: Internal VM1



In this the Internal machine i.e. VM1 (192.168.130.135) is generating ping messages and is continuously sending it to the external machine i.e. VM3 (192.168.140.188). This ping is sending 200 packets and they are sent periodically at a delay of 20ms from each other.

vm3@vm3-Standard-PC-Q35-ICH9-2009:-\$ ifconfig			
enpis0: fl ags=4103<up,broadcas< del="">T,RUNNING,MULTICAST> mtu 1500</up,broadcas<>			
inet 192.168.140.188 netmask 255.255.255.0 broadcast 192.168.140.255			
i det6_fe80::6088:96b1/ 24dd:fa25_prefixlen_64_scopeid_0x20 <link/>			
ether 52:54:00:d6:10:87 txqueuelen 1000 (Ethernet)			
RX packets 26886 bytes 2404426 (2.4 MB)			
RX errors 0 dropped 17434 overruns 0 frame 0			
TX packets 7951 bytes 737328 (737.3 KB)			
TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0			
· 전· 이상 ·			
lo: flags=73 <up,loopback,running> mtu 65536</up,loopback,running>			
inet 127.0.0.1 netmask 255.0.0.0			
inet6 ::1 prefixlen 128 scopeid 0x10 <host></host>			
loop txqueuelen 1000 (Local Loopback)			
RX packets 804 bytes 74284 (74.2 KB)			
RX errors 0 dropped 0 overruns 0 frame 0			
TX packets 804 bytes 74284 (74.2 KB)			
TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0			
vm3@vm3-Standard-PC-Q35-ICH9-2009:-\$			

Fig: IP address of external machine VM3



Fig: Command for DoS attack prevention in VM2 (Firewall)

What this command does is it enables the DoS prevention mechanism and if the number of packets from an IP address is more than 100 then the Firewall blocks the IP address.



Fig: Rules description

As we can see in the Fig above the rule with rule_id: 100, tells to allow the packets with which have destination address as **192.168.140.188**



[Status]: Dropped [PPT] : 4.311e-05
<pre>[Ethernet][IPv4];1m[TCP] [Src MAC]: 52:54:00:f7:69:35, [Dstn MAC]: 52:54:00:7f:ab:9d [Src IP]: 192.168.130.135, [Dstn IP]: 192.168.140.188 [Status]: Allowed [PPT] : 7.084e-05</pre>
<pre>[Ethernet][IPv4];1m[TCP] [Src MAC]: 52:54:00:41:10:ec</pre>

Fig: Firewall ON, DoS is not detected yet

Since the number of Packets is not more than 100 DoS is not detected yet.

[Status]: Dropped [PPT] : 0.00013958
DoS Detected [Ethernet][IPv4]:1m[TCP] [Src MAC]: 52:54:00:f7:69:35, [Dstn MAC]: 52:54:00:7f:ab:9d [SrcIP]: 192.168.130.135, [Dstn IP]: 192.168.140.188 [Status]: Dropped [PPT] : 0.00010329
DoS Detected [Ethernet][IPv4] <mark>:1m[TCP]</mark> [Src MAC]: 52:54:00:f7:69:35, [Dstn MAC]: 52:54:00:7f:ab:9d

Fig: Firewall ON, DoS is Detected

In the above Figure the number of packets has crossed 100 mark hence the Firewall has started to block the packets.

In the figure below, we can see that once the number of packets reaches 100 then the IP address of VM1 (192.168.130.135) is blocked and no more packets are incoming from it. Thus the number of allowed packets plateaus after a point.





Fig: Cumulative graph of allowed packets and Dropped packets

This figure clearly demonstrates how the packets gets dropped when a firewall detects DDoS.